# Face numbers: the upper bound side of the story

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#### I. Basics on Polytopes

A polytope is the convex hull of *finitely many* points in  $\mathbb{R}^d$ Equivalently, it is a bounded intersection of *finitely many* closed half-spaces in  $\mathbb{R}^d$ **Example:** a simplex is the convex hull of affinely independent points Polytopes are studied and have applications in combinatorics, discrete geometry, optimization, analysis, statistics, ...



#### Faces and face numbers

- The dimension of a polytope *P* is the dimension of its affine hull
- A face of *P* is the intersection of *P* with a supporting hyperplane
- A face F of P is itself a polytope; it is an *i*-face if dim F = i

Given a *d*-polytope we can count how many vertices, edges, 2-faces,..., (d-1)-faces (also known as facets) it has:

#### Motivation --- the Upper Bound Problem

- What is the largest number of *i*-faces that a *d*-polytope with *n* vertices can have?
- Connection to optimization: the dual form of this question is "What is the largest number of vertices that a d-polytope defined by n linear constraints can have?"

To start, we need to look for polytopes with many faces This leads us to the cyclic polytopes (discovered and rediscovered by Carathéodory, Gale, Motzkin,...)

#### Cyclic polytopes



The cyclic polytope, C(d, n), is defined as  $conv(M(t_1), M(t_2), ..., M(t_n))$ 

## Properties of cyclic polytopes

- C(d, n) is a d-dimensional simplicial polytope on n vertices (i.e., all facets are simplices)
- The combinatorial type of C(d, n) is independent of the choice of t<sub>1</sub>, t<sub>2</sub>, ..., t<sub>n</sub>: there is a complete characterization of the vertex sets of facets due to Gale (Gale's evenness condition)
- C(d, n) is  $\lfloor \frac{d}{2} \rfloor$ -neighborly: every set of  $\leq \lfloor \frac{d}{2} \rfloor$  vertices forms the vertex set of a face. Thus

$$f_{i-1}(C(d,n)) = \binom{n}{i} \quad \forall i \le \lfloor \frac{d}{2} \rfloor$$

#### A digression: simplicial complexes and simplicial spheres

Def  $\Delta \subseteq 2^V$  is a simplicial complex on a finite vertex set V if

- $\{v\} \in \Delta \quad \forall v \in V$
- $F \in \Delta$ ,  $G \subset F \Rightarrow G \in \Delta$

Elements of V are vertices, elements of  $\Delta$  are faces

A face F is an *i*-face if |F| = i + 1; the number of *i*-faces is  $f_i(\Delta)$ 

Simplicial complex  $\Delta \rightarrow$  Topological space  $\|\Delta\|$  = the geometric realization of  $\Delta$ 

 $\Delta$  is a simplicial (d - 1)-sphere if  $\|\Delta\|$  is homeomorphic to  $S^{d-1}$  $\Delta$  is a simplicial manifold if  $\|\Delta\|$  is homeomorphic to a (closed) manifold

## The Upper Bound Theorems

#### The Upper Bound Conjecture

Motzkin, 1957: Among all *d*-polytopes with *n* vertices, the cyclic polytope simultaneously maximizes all the face numbers  $f_i(P) \le f_i(C(d, n)) \forall i$ 

V. Klee, 1964: Among all Eulerian complexes of dimension d - 1 with n vertices, the boundary complex of the cyclic polytope simultaneously maximizes all the face numbers

Eulerian complexes include all simplicial spheres, all odddim manifolds, all even-dim manifolds with Euler char 2

#### **The Upper Bound Theorems:**

≻ P. McMullen, 1970: The UBC holds for all polytopes

- **R. Stanley, 1975:** The UBC holds for all simplicial spheres
- ≻N, 1998: The UBC holds for all Eulerian manifolds

#### Comments

McMullen's proof uses *shellability* of polytopes

Stanley's proof relies on the theory of Stanley-Reisner rings and specifically on the properties of Cohen-Macaulay rings

The proof for manifolds relies on the properties of **Buchsbaum rings** 

C(d, n) in the statement of the UBT can be replaced with any  $\lfloor \frac{d}{2} \rfloor$ -neighborly d-polytope or (d - 1)-sphere with n vertices. This leads us to the question of how many  $\lfloor \frac{d}{2} \rfloor$ -neighborly (d - 1)-spheres with n vertices are there?

#### II. There are many more spheres than polytopes

- Let c(d, n) = # of simplicial d-polytopes with n (labeled) vertices
- Let s(d, n) = # of simplicial (d 1)-spheres with n (labeled) vertices

Steinitz's theorem implies that c(3,n) = s(3,n) **Theorem (Goodman-Pollack; Alon 1986)** For  $d \ge 4$ ,  $c(d,n) = 2^{\Theta(n \log n)}$ , i.e.,  $2^{a_d n \log n} \le c(d,n) \le 2^{A_d n \log n}$  for some constants  $a_d$ ,  $A_d > 0$ .

**Theorem** (Kalai, 1988; Pfeifle-Ziegler, 2004; Nevo-Santos-Wilson, 2016) For  $d \ge 4$ ,  $2^{\Omega(n^{\lfloor \frac{d}{2} \rfloor})} \le s(d, n) \le 2^{O(n^{\lfloor \frac{d}{2} \rfloor} \log n)}$  (e.g.,  $2^{\Omega(n^2)} \le s(4, n) \le 2^{O(n^2 \log n)}$ )

## What proportion of d-polytopes are $\lfloor \frac{d}{2} \rfloor$ -neighborly ?

Results of Shemer and Padrol indicate that "most of d-polytopes are  $\lfloor \frac{d}{2} \rfloor$ -neighborly as  $n \to \infty$ ":

#### Theorem (Shemer, 1982; Padrol, 2013)

- There are  $2^{\Theta(n \log n)} \lfloor \frac{d}{2} \rfloor$ -neighborly simplicial d-polytopes with n vertices. Padrol's lower bound on the number of  $\lfloor \frac{d}{2} \rfloor$ -neighborly d-polytopes with n vertices is currently the best known lower bound on c(d, n)
- There are at least  $2^{\Omega(n \log n)} \lfloor \frac{d}{2} \rfloor$ -neighborly (d 1)-spheres with n vertices arising from non-realizable oriented matroids

## What proportion of (d - 1)-spheres are $\lfloor \frac{d}{2} \rfloor$ -neighborly ?

Let sn(d, n) = # of  $\lfloor \frac{d}{2} \rfloor$ -neighborly simplicial (d - 1)-spheres with n vertices

**Conjecture** (Kalai, 1988) For all  $d \ge 4$ ,  $\lim_{n \to \infty} \frac{\log sn(d,n)}{\log s(d,n)} = 1$ 

**Theorem** (N-Zheng, 2021+) There are many neighborly spheres: for all  $d \ge 5$ ,  $\operatorname{sn}(d,n) \ge 2^{\Omega(n \lfloor \frac{d-1}{2} \rfloor)}$ 

The proof is by construction based on Kalai's squeezed balls --- certain subcomplexes of the boundary complex of C(d,n)

[For comparison, recall that  $2^{\Omega(n^{\lfloor \frac{d}{2} \rfloor})} \le s(d, n) \le 2^{\Omega(n^{\lfloor \frac{d}{2} \rfloor} \log n)}$ ]

## Summary so far

- For  $d \ge 4$ , there are many more simplicial (d 1)-spheres than d-polytopes
- Nonetheless, d-polytopes with n vertices and (d 1)-spheres with n vertices (and even Eulerian (d - 1)-manifolds with n vertices) satisfy the same Upper Bound Theorem

In fact, recent very exciting news (due to Adiprasito, and Papadakis and Petrotou) is that the set of f-vectors of simplicial (d - 1)-spheres coincides with the set of f-vectors of simplicial d-polytopes

- The maximizers are given by  $\lfloor \frac{d}{2} \rfloor$ -neighborly d-polytopes and  $\lfloor \frac{d}{2} \rfloor$ -neighborly (d 1)-spheres
- There are many  $\lfloor \frac{d}{2} \rfloor$ -neighborly d-polytopes; there are also many  $\lfloor \frac{d}{2} \rfloor$ -neighborly (d-1)-spheres

#### III. Cs polytopes and cs spheres

- A polytope  $P \subset \mathbb{R}^d$  is centrally symmetric if  $x \in P \Leftrightarrow -x \in P$
- A simplicial sphere  $\Delta$  is centrally symmetric if  $\exists \varphi: V \to V$  such that  $\varphi(F) \in \Delta$ ,  $\varphi(\varphi(F)) = F$ , but  $\varphi(F) \neq F \quad \forall \emptyset \neq F \in \Delta$ (The vertices v and  $\varphi(v)$  are called antipodal)

Note: if  $\Delta$  is cs, then v and  $\varphi(v)$  are **not** connected by an edge!



#### The Upper Bound Problem for cs polytopes and spheres

#### **Problems:**

- What restrictions does being cs impose on the *f*-vectors?
- More specifically, what is the largest number of *i*-faces that a cs *d*-polytope with *n* vertices can have? What is the largest number of *i*-faces that a cs (d 1)-sphere with *n* vertices can have?

#### **Motivation**:

Donoho, and Rudelson and Vershynin observed that cs polytopes with many faces have applications in sparse signal reconstruction and coding theory

#### Cs neighborliness

A cs simplicial sphere  $\Delta$  is cs-k-neighborly if every set of  $\leq k$  vertices of  $\Delta$  no two of which are antipodal is the vertex set of a face of  $\Delta$ 

Examples:

1) the *d*-dimensional cross-polytope  $C_d^* \coloneqq \operatorname{conv}(\pm e_1, \pm e_2, \dots, \pm e_d)$  is cs-*d*-neighborly



2) McMullen-Shephard, 1968: conv $(\pm e_1, \pm e_2, \dots, \pm e_d, \pm (e_1 + \dots + e_d))$  is cs- $\lfloor \frac{d}{2} \rfloor$ -neighborly

#### How neighborly can a cs polytope be?

Keeping the cyclic polytope in mind, we might expect the answer to be  $\left[\frac{a}{2}\right]$ . However:

**Theorem** (McMullen-Shephard, 1968; d = 4 case is due to Grünbaum, 1967) A cs d-polytope with  $\geq 2(d + 2)$  vertices cannot be cs- $\left(\left\lfloor \frac{d+1}{3} \right\rfloor + 1\right)$ -neighborly (e.g., a cs 4-polytope with 12 vertices cannot be cs-2-neighborly)

**Theorem (Burton, 1991)** A cs d-polytope with a sufficiently large number (about  $d^d$ ) of vertices cannot be even cs-2-neighborly

## Cs-2-neighborliness of cs polytopes

**Theorem** Let  $d \ge 3$  be any integer.

1) [Linial-N, 2006] A cs d-polytope with  $2^d$  or more vertices cannot be even cs-2-neighborly

The proof is based on the volume argument going back to Danzer-Grünbaum

2) [N, 2018] There exists a cs d-polytope with  $2^{d-1} + 2$  vertices that is cs-2-neighborly.

Embed the (d-1)-cube  $C_{d-1}$  in  $\mathbb{R}^d$  as  $[-1,1]^{d-1} \times \{0\}$  and perturb its vertices using the d-th dimension

Thus, the maximum number of vertices that a cs-2-neighborly d-polytope can have lies in  $[2^{d-1} + 2, 2^d - 2]$ .

**Open**: What is this number? [For d = 3, 4, it is  $2^{d-1} + 2$ .]

#### From non-neighborliness to *f*-numbers

What is the value of  $fmax(d, n; 1) \coloneqq max\{f_1(P) : P \text{ is a cs polytope, dim } P = d, f_0(P) = n\}$ ? By non cs-2-neighborliness, if  $n \ge 2^d$ , then  $fmax(d, n; 1) < {n \choose 2} - \frac{n}{2}$ 

**Theorem** [Barvinok-N, 2008; Barvinok-Lee-N, 2013] For an even  $d \ge 4$ ,

$$\left(1 - 3 \cdot \left(\sqrt{3}\right)^{-d}\right) \binom{n}{2} \le \operatorname{fmax}(d, n; 1) \le \left(1 - 2^{-d}\right) \frac{n^2}{2}$$

For 
$$d = 4$$
,  $\frac{3}{4} \cdot \frac{n^2}{2} - O(n) \le \text{fmax}(4, n; 1) \le \frac{15}{16} \cdot \frac{n^2}{2}$ 

Wide open: what is the value of fmax(4, n; 1)?

## IV. cs neighborliness of cs spheres

We saw that cs-neighborliness of cs polytopes is quite restricted What about *cs simplicial spheres*?

**Theorem (Adin, 1991; Stanley)** Among all cs simplicial (d - 1)-spheres on n vertices, a cs- $\lfloor \frac{d}{2} \rfloor$ -neighborly simplicial sphere simultaneously maximizes all the face numbers, assuming such a sphere exists

#### Does it exist?

Grünbaum, late 60s: there is a cs simplicial 3-sphere with 12 vertices that is cs-2-neighborly.

Jockusch, 1995: For every  $m \ge 4$ , there exists a cs simplicial 3-sphere with 2m vertices that is cs-2-neighborly.

Lutz, 1999: there is a cs simplicial 5-sphere with 16 vertices that is cs-3-neighborly; there is also a cs simplicial 7-sphere with 18 vertices that is cs-4-neighborly

(In contrast, by [McMullen-Shephard, 1968] there are no such cs polytopes)

#### It does exist!

#### Theorem (N-Zheng, 2020)

For every  $d \ge 4$  and  $m \ge d$ , there is a cs (d - 1)-sphere with 2m vertices,  $\Delta_m^{d-1}$ , that is cs- $\lfloor \frac{d}{2} \rfloor$ -neighborly

(In fact, for  $m \gg d$ , there are at least two non-isomorphic constructions)

This together with Adin's and Stanley's work leads to

The Upper Bound Theorem for cs spheres Among all cs simplicial (d - 1)-spheres on 2m vertices,  $\Delta_m^{d-1}$  simultaneously maximizes all the face numbers

## Summary and Open problems

The *f*-vectors of simplicial spheres/polytopes without symmetry satisfy the same UBT.

The situation for cs spheres and cs polytopes is drastically different: the Upper Bound Problem for cs simplicial spheres is now completely resolved, but for cs *polytopes*, there is not even a plausible Upper Bound Conjecture

In fact, we don't even know

- > What is the maximum possible number of edges that a cs 4-polytope with 2m vertices can have?
- ➢What is the maximum possible number of vertices that a cs-2-neighborly *d*-polytope can have?

There are many more remaining mysteries, but let me stop here

## THANK YOU!